Critical independent sets and König-Egerváry graphs

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Abstract

A set S of vertices is *independent* in a graph G, and we write $S \in \text{Ind}(G)$, if no two vertices from S are adjacent, and $\alpha(G)$ is the cardinality of an independent set of maximum size, while core(G) denotes the intersection of all maximum independent sets [17].

G is called a König–Egerváry graph if its order equals $\alpha(G) + \mu(G)$, where $\mu(G)$ denotes the size of a maximum matching. The number $def(G) = |V(G)| - 2\mu(G)$ is the deficiency of G [21].

The number $d(G) = \max\{|S| - |N(S)| : S \in \operatorname{Ind}(G)\}$ is the *critical difference* of G. An independent set A is *critical* if |A| - |N(A)| = d(G), where N(S) is the neighborhood of S, and $\alpha_c(G)$ denotes the maximum size of a critical independent set [26].

In [14] it was shown that G is König–Egerváry graph if and only if there exists a maximum independent set that is also critical, i.e., $\alpha_c(G) = \alpha(G)$.

In this paper we prove that:

- (i) $d(G) = |\text{core}(G)| |N(\text{core}(G))| = \alpha(G) \mu(G) = def(G)$ hold for every König–Egerváry graph G;
- (ii) G is König–Egerváry graph if and only if each maximum independent set of G is critical.

Keywords: independent set, maximum matching, critical difference, critical independent set, deficiency, core.

1 Introduction

Throughout this paper G = (V, E) is a finite, undirected, loopless and without multiple edges graph with vertex set V = V(G) and edge set E = E(G). If $X \subset V$, then G[X] is the subgraph of G spanned by X. By G - W we mean the subgraph G[V - W], if $W \subset V(G)$. For $F \subset E(G)$, by G - F we denote the partial subgraph of G obtained by deleting the edges of F, and we use G - e, if $W = \{e\}$. The neighborhood of a vertex $v \in V$ is the set $N(v) = \{w : w \in V \text{ and } vw \in E\}$, while $N(A) = \bigcup \{N(v) : v \in A\}$ and $N[A] = A \cup N(A)$ for $A \subset V$.

A set $S \subseteq V(G)$ is independent if no two vertices from S are adjacent, and by $\operatorname{Ind}(G)$ we mean the set of all the independent sets of G. An independent set of maximum size will be referred to as a maximum independent set of G, and the independence number of G is $\alpha(G) = \max\{|S| : S \in \operatorname{Ind}(G)\}$.

Let us denote the set $\{S : S \text{ is a maximum independent set of } G\}$ by $\Omega(G)$, and let $\operatorname{core}(G) = \cap \{S : S \in \Omega(G)\}$ [17]. A set $A \subseteq V(G)$ is a local maximum independent set of G if $A \in \Omega(G[N[A]])$ [16].

Theorem 1.1 [22] Every local maximum independent set of a graph is a subset of a maximum independent set.

A matching (i.e., a set of non-incident edges of G) of maximum cardinality $\mu(G)$ is a maximum matching, and a perfect matching is one covering all vertices of G.

It is well-known that

$$||V|/2| + 1 \le \alpha(G) + \mu(G) \le |V|$$

hold for any graph G = (V, E). If $\alpha(G) + \mu(G) = |V|$, then G is called a König-Egerváry graph. We attribute this definition to Deming [6], and Sterboul [25]. These graphs were studied in [3, 11, 15, 18, 19, 20, 21, 24], and generalized in [2, 23].

According to a well-known result of König [10], and Egerváry [8], any bipartite graph is a König-Egerváry graph. This class includes non-bipartite graphs as well (see, for instance, the graphs H_1 and H_2 in Figure 1).

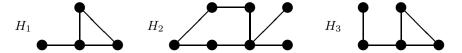


Figure 1: Only H_3 is not a König-Egerváry graph, as $\alpha(H_3) + \mu(H_3) = 4 < 5 = |V(H_3)|$.

It is easy to see that if G is a König-Egerváry graph, then $\alpha(G) \ge \mu(G)$, and that a graph G having a perfect matching is a König-Egerváry graph if and only if $\alpha(G) = \mu(G)$.

The number $d(G) = \max\{|S| - |N(S)| : S \in \text{Ind}(G)\}$ is called the *critical difference* of G. An independent set A is *critical* if |A| - |N(A)| = d(G), and the *critical independence* number $\alpha_c(G)$ is the cardinality of a maximum critical independent set [26]. Clearly, $\alpha_c(G) \leq \alpha(G)$ holds for any graph G. It is known that the problem of finding a critical independent set is polynomially solvable [1, 26].

Proposition 1.2 [13] If S is a critical independent set, then there is a matching from N(S) into S.

If S is an independent set of a graph G and H = G - S, then we write G = S * H. Evidently, any graph admits such representations. For instance, if $E(H) = \emptyset$, then G = S * H is bipartite; if H is complete, then G = S * H is a split graph [9].

Proposition 1.3 [18] G is a König-Egerváry graph if and only if $G = H_1 * H_2$, where $V(H_1) \in \Omega(G)$ and $|V(H_1)| \ge \mu(G) = |V(H_2)|$.

Let M be a maximum matching of a graph G. To adopt Edmonds's terminology [7], we recall the following terms for G relative to M. An alternating path from a vertex x to a vertex y is a x,y-path whose edges are alternating in and not in M. A vertex x is exposed relative to M if x is not the endpoint of a heavy edge. An odd cycle C with $V(C) = \{x_0, x_1, ..., x_{2k}\}$ and $E(C) = \{x_i x_{i+1} : 0 \le i \le 2k-1\} \cup \{x_{2k}, x_0\}$, such that $x_1 x_2, x_3 x_4, ..., x_{2k-1} x_{2k} \in M$ is a blossom relative to M. The vertex x_0 is the base of the blossom. The stem is an even length alternating path joining the base of a blossom and an exposed vertex for M. The base is the only common vertex to the blossom and the stem. A flower is a blossom and its stem. A posy consists of two (not necessarily disjoint) blossoms joined by an odd length alternating path whose first and last edges belong to M. The endpoints of the path are exactly the bases of the two blossoms. The following result of Sterboul, characterizes König-Egerváry graphs in terms of forbidden configurations.

Theorem 1.4 [25] For a graph G, the following properties are equivalent:

- (i) G is a König-Egerváry graph;
- (ii) there exist no flower and no posy relative to some maximum matching M;
- (iii) there exist no flower and no posy relative to any maximum matching M.

In [20] is given a characterization of König-Egerváry graphs having a perfect matching, in terms of certain forbidden subgraphs with respect to a specific perfect matching of the graph. In [12] is given the following characterization of König-Egerváry graphs in terms of excluded structures.

Theorem 1.5 [12] Let M be a maximum matching in a graph G. Then G is a König-Egerváry graph if and only if G does not contain one of the forbidden configurations, depicted in Figure 2, with respect to M.

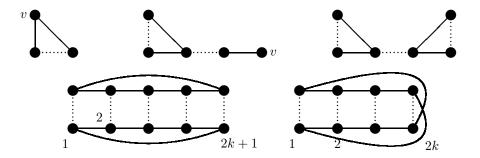


Figure 2: Forbidden configurations. The vertex v is not adjacent to the matching edges (namely, dashed edges).

In [14] it was shown that G is a König-Egerváry graph if and only if $\alpha_c(G) = \alpha(G)$, thus giving a positive answer to the Graffiti.pc 329 conjecture [5].

The deficiency of G, denoted by def(G), is defined as the number of exposed vertices relative to a maximum matching [21]. In other words, $def(G) = |V(G)| - 2\mu(G)$.

In this paper we prove that the critical difference for a König-Egerváry graph G is given by

$$d(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))| = \alpha(G) - \mu(G) = \operatorname{def}(G),$$

and using this finding, we show that G is a König-Egerváry graph if and only if each of its maximum independent sets is critical.

2 Results

Proposition 2.1 Every critical independent set is a local maximum independent set.

Proof. Suppose, on the contrary, that there is a critical independent set S such that $S \notin \Psi(G)$, i.e., there exists some independent set $A \subseteq N[S]$, larger than S. It follows that $|A \cap N(S)| > |S - S \cap A|$, and this contradicts the fact that, according to Proposition 1.2, there is a matching from $A \cap N(S)$ to S, in fact, from $A \cap N(S)$ to $S - S \cap A$.

The converse of Proposition 2.1 is not true; e.g., the set $\{d, h\}$ is a local maximum independent set of the graph G_1 from Figure 3, but it is not critical.

Using Theorem 1.1, we easily deduce the following result.

Corollary 2.2 [4] Every critical independent set is contained in some maximum independent set.

Theorem 2.3 If G is a König-Egerváry graph, then

- (i) $[18] N(\operatorname{core}(G)) = \bigcap \{V(G) S : S \in \Omega(G)\};$
- (ii) $[19] \alpha(G) + |\cap \{V(G) S : S \in \Omega(G)\}| = \mu(G) + |\cap \{S : S \in \Omega(G)\}|;$
- (iii) /19/G N[core(G)] has a perfect matching and it is also a König-Egerváry graph.

Let us notice that for non-König-Egerváry graphs every relation between $\alpha(G) - \mu(G)$ and $|\operatorname{core}(G)| - |N(\operatorname{core}(G))|$ is possible.

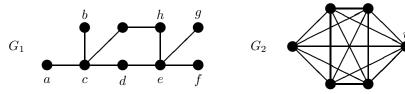


Figure 3: $\alpha(G_1) = 6$, $\mu(G_1) = 3$, $\operatorname{core}(G_1) = \{a, b, d, g, f\}$ and $N(\operatorname{core}(G_1)) = \{c, e\}$, while $\alpha(G_2) = 4$, $\mu(G_2) = 3$, $\operatorname{core}(G_2) = \{x, y, z\}$, and $N(\operatorname{core}(G_2)) = \{v\}$.

The non-König-Egerváry graphs from Figure 3 satisfy:

$$\alpha(G_1) - \mu(G_1) = 3 = |\operatorname{core}(G_1)| - |N(\operatorname{core}(G_1))|$$

and

$$\alpha(G_2) - \mu(G_2) = 1 < 2 = |\operatorname{core}(G_2)| - |N(\operatorname{core}(G_2))|.$$

The opposite direction of the above inequality may be found in $G_3 = K_{2n} - e, n \ge 3$:

$$\alpha(G_3) - \mu(G_3) = 2 - n > 4 - 2n = 2 - (2n - 2) = |\operatorname{core}(G_3)| - |N(\operatorname{core}(G_3))|.$$

Theorem 2.4 If G is König-Egerváry graph, then the following equalities hold

$$d(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))| = \alpha(G) - \mu(G) = \operatorname{def}(G).$$

Proof. Firstly, let us prove that $\alpha(G) - \mu(G) \ge |S| - |N(S)|$ holds for every $S \in \text{Ind}(G)$, i.e., $d(G) \le \alpha(G) - \mu(G)$. If $\alpha(G) = \mu(G)$, then G has a perfect matching and

$$|S| - |N(S)| < 0 = \alpha(G) - \mu(G)$$

holds for every $S \in \text{Ind}(G)$.

Suppose that $\alpha(G) > \mu(G)$. Let $S_0 \in \Omega(G)$ and M be a maximum matching, i.e., $|M| = |V(G) - S_0| = \mu(G)$. Assume that $S \in \operatorname{Ind}(G)$ satisfies |S| - |N(S)| > 0. Then one can write $S = S_1 \cup S_2 \cup S_3$, where $S_3 \subseteq V(G) - S_0$, $S_1 \cup S_2 \subset S_0$, $S_1 \cap S_2 = \emptyset$, and S_2 contains every $v \in S$ matched by M with some vertex of $V(G) - S_0$. Since M is a maximum matching, we obtain that $|S_2| - |N(S_2)| \le 0$ and $|S_3| - |N(S_3)| \le 0$. Consequently, we infer that

$$\alpha(G) - \mu(G) = |S_0| - |V(G) - S_0| \ge |S_1| \ge |S| - |N(S)|,$$

as required (see Figure 4 for various examples of S).

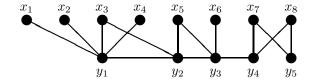


Figure 4: $S_0 = \{x_i : 1 \le i \le 8\}, M = \{y_1x_4, y_2x_5, y_3x_6, y_4x_7, y_5x_8\}, S = S_1 \cup S_2 \cup S_3,$ where $S_2 = \{x_5\}, S_3 = \{y_4, y_5\},$ while S_1 belongs to $\{\{x_1, x_2\}, \{x_1x_3\}, \{x_3\}\}.$

The fact that core(G) is an independent set of G ensures that

$$\alpha(G) - \mu(G) \ge |\operatorname{core}(G)| - |N(\operatorname{core}(G))|$$
.

Since G is a König-Egerváry graph, we get that

$$\alpha(G) + \mu(G) = |V(G)| = |\operatorname{core}(G)| + |N(\operatorname{core}(G))| + |V(G - N[\operatorname{core}(G)])|$$
.

Assuming that

$$\alpha(G) - \mu(G) > |\operatorname{core}(G)| - |N(\operatorname{core}(G))|,$$

we obtain the following contradiction

$$\begin{aligned} 2\alpha(G) &> 2\left|\operatorname{core}(G)\right| + \left|V(G - N[\operatorname{core}(G)])\right| \\ &= 2\left|\operatorname{core}(G)\right| + 2\alpha(G - N[\operatorname{core}(G)]) = 2\alpha(G), \end{aligned}$$

because $|V(G - N[\operatorname{core}(G)])| = 2\alpha (G - N[\operatorname{core}(G)])$ by Theorem 2.3(iii).

Therefore, we get that $\alpha(G) - \mu(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))|$. Actually, this equality immediately follows from Theorem 2.3(i), (ii), but the current way of proof exploits different aspects of $\operatorname{Ind}(G)$.

Further, using the inequality $d(G) \leq \alpha(G) - \mu(G)$ and the equality

$$\alpha(G) - \mu(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))|,$$

we finally deduce that

$$|\operatorname{core}(G)| - |N(\operatorname{core}(G))| \le \max\{|S| - |N(S)| : S \in \operatorname{Ind}(G)\} = d(G)$$

$$\le \alpha(G) - \mu(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))|,$$

i.e.,

$$\alpha(G) - \mu(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))| = d(G).$$

Since G is a König-Egerváry graph, we infer that

$$\alpha(G) - \mu(G) = \alpha(G) + \mu(G) - 2\mu(G) = |V(G)| - 2\mu(G) = def(G),$$

and this completes the proof.

Corollary 2.5 If G is a König-Egerváry graph, then d(G) = 0 if and only if G has a perfect matching.

Remark 2.6 There exist non-König-Egerváry graphs enjoying the equalities

$$d(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))| = \alpha(G) - \mu(G),$$

see, for instance, the graph G from Figure 5.

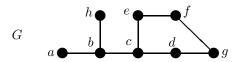


Figure 5: *G* has $\alpha(G) = 4$, $\mu(G) = 3$, $\text{core}(G) = \{a, h\}$ and $N(\text{core}(G)) = \{b\}$.

Theorem 2.7 The following assertions are equivalent:

- (i) G is a König-Egerváry graph;
- (ii) there is $S \in \Omega(G)$, such that S is critical, i.e., $\alpha_c(G) = \alpha(G)$;
- (iii) every $S \in \Omega(G)$ is critical.

Proof. (i) \Longrightarrow (iii) Let $S \in \Omega(G)$, $A = S - \operatorname{core}(G)$ and $B = V(G) - S - N(\operatorname{core}(G))$. By Theorem 2.3(iii), we infer that |A| = |B|, since $G - N[\operatorname{core}(G)]$ has a perfect matching. Hence, we obtain that

$$|S| - |N(S)| = |A| + |\operatorname{core}(G)| - (|B| + |N(\operatorname{core}(G))|$$

= $|\operatorname{core}(G)| - |N(\operatorname{core}(G))|$.

In other words, according to Theorem 2.4, the equality |S| - |N(S)| = d(G) is true for every $S \in \Omega(G)$.

- $(iii) \Longrightarrow (ii)$ It is clear.
- $(ii) \Longrightarrow (i)$ This was done in [14]. For the sake of completeness we add the proof.

There is a critical independent set S with $|S| = \alpha_c(G) = \alpha(G)$. By Proposition 1.2, there exists a matching M from N(S) into S, and clearly, $|M| = |N(S)| = \mu(G)$. Hence, we finally obtain that $|V(G)| = |S| + |N(S)| = \alpha(G) + \mu(G)$, i.e., G is a König-Egerváry graph. \blacksquare

3 Conclusions

In this paper we give a new characterization of König-Egerváry graphs. On the one hand, it is similar in form to Sterboul's theorem [25]. On the other hand it extends Larson's finding [14]. We found that the critical difference of a König-Egerváry graph G is given by

$$d(G) = |\operatorname{core}(G)| - |N(\operatorname{core}(G))| = \alpha(G) - \mu(G) = \operatorname{def}(G).$$

It seems interesting to find other families of graphs satisfying these equalities.

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